# Time-Uniform Confidence Sequences from Universal Gambling 

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- For reliable inference, we need to quantify confidence


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which is equivalent to

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\mu \in\left(\hat{\mu}_{t}-\sqrt{\frac{1}{t} \frac{\operatorname{Var}(Y)}{\delta}}, \hat{\mu}_{t}+\sqrt{\frac{1}{t} \frac{\operatorname{Var}(Y)}{\delta}}\right) \text { with prob. } \geq 1-\delta
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## Confidence Sets

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- Now, suppose we wish to decide to keep or stop sampling $Y_{t}$ to estimate $\mu$ given confidence level on the fly (sequentially)
- For such online data processing, we need to construct a sequence of confidence intervals that is valid at any time


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- Originally studied by Darling and Robbins (1967); Lai (1976), and recently resurrected by some statisticians (Ramdas et al., 2020; Waudby-Smith and Ramdas, 2020a,b; Howard et al., 2021) and computer scientists (Jun and Orabona, 2019; Orabona and Jun, 2021)


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## Ville's inequality (Ville, 1939)

For a nonnegative supermartingale sequence $\left(W_{t}\right)_{t=0}^{\infty}$ with $W_{0}>0$,

$$
\mathrm{P}\left\{\sup _{t \geq 1} \frac{W_{t}}{W_{0}} \geq \frac{1}{\delta}\right\} \leq \delta
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- Wait, what is gambling?
- As a slight detour, let's review canonical gambling problems and some universal gambling strategies


## Universal Gambling

## Coin Betting

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- The recursive equation:


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\mathbf{W}_{t}=\mathbf{W}_{t-1} 2 q_{t}^{\mathbb{1}\left\{y_{t}=1\right\}}\left(1-q_{t}\right)^{\mathbb{1}\left\{y_{t}=0\right\}}=\mathbf{W}_{t-1} 2 q\left(y_{t} \mid y^{t-1}\right)
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$$

- Cumulative wealth: starting with $\$ \mathrm{~W}_{0}$,

$$
\mathbf{W}_{T}=\mathbf{W}_{0} \prod_{t=1}^{T} 2 q\left(y_{t} \mid y^{t-1}\right)=\mathbf{W}_{0} 2^{T} q\left(y^{T}\right)
$$

where $q\left(y^{T}\right):=\prod_{t=1}^{T} q\left(y_{t} \mid y^{t-1}\right)$

## Universality and Minimax Optimality

- Let $\mathrm{W}_{t}:=\mathrm{W}^{q}\left(y^{t}\right)$ for a betting strategy $\left(q\left(\cdot \mid y^{t-1}\right)\right)_{t=1}^{\infty}$
- For some $\mathcal{P}=\{$ reference strategies $p\}$, track the best performance of $\mathcal{P}$ in hindsight


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- Wealth ratio

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- Regret $=\log$ (wealth ratio)

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- Regret w.r.t. the best reference strategy

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- Worst-case regret w.r.t. the best reference strategy

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- The best strategy is called minimax optimal

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## Coin Betting $\equiv$ Probability Assignment

- Note: $\frac{\mathrm{W}^{p}\left(y^{T}\right)}{\mathrm{W}^{q}\left(y^{T}\right)}=\frac{\mathrm{W}_{0} 2^{T} p\left(y^{T}\right)}{\mathrm{W}_{0} 2^{T} q\left(y^{T}\right)}=\frac{p\left(y^{T}\right)}{q\left(y^{T}\right)}$ by definition


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- The cumulative regret w.r.t. a reference probability $p\left(y^{t}\right)$ is

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$\therefore$ universal compression $\rightarrow$ universal betting!

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- Fact: $p_{\theta^{*}}$ is optimal if $y^{T} \sim$ i.i.d. $\operatorname{Bern}\left(\theta^{*}\right)$ (a.k.a. Kelly betting)


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\mathrm{W}^{\theta}\left(y^{T}\right):=\mathrm{W}_{0} 2^{T} p_{\theta}\left(y^{T}\right)
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where $p_{\theta}\left(y^{T}\right)$ is the "probability" under $y^{T} \sim$ i.i.d. $\operatorname{Bern}(\theta)$

- Fact: $p_{\theta^{*}}$ is optimal if $y^{T} \sim$ i.i.d. $\operatorname{Bern}\left(\theta^{*}\right)$ (a.k.a. Kelly betting)
- Krichevsky-Trofimov (KT) probability assignment (Krichevsky and Trofimov, 1981)

$$
q_{\mathrm{KT}}\left(1 \mid y^{t-1}\right):=\frac{1}{t}\left(\sum_{i=1}^{t-1} y_{i}+\frac{1}{2}\right)
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## Example: Constant Bettors

- $\mathcal{P}=\left\{p_{\theta}(\cdot): \theta \in[0,1]\right\}$, where $p_{\theta}\left(1 \mid y^{t-1}\right)=\theta$
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- Asymptotically minimax optimal (Xie and Barron, 2000)

$$
\max _{\theta \in[0,1]} \max _{y^{T}} \log \frac{p_{\theta}\left(y^{T}\right)}{q_{\mathrm{KT}}\left(y^{T}\right)}=\frac{1}{2} \log T+\frac{1}{2} \log \frac{\pi}{2}+o(1)
$$

## Mixture Probability

- The KT probability $q_{\mathrm{KT}}\left(\cdot \mid y^{t-1}\right)$ is induced by a mixture probability, i.e.,

$$
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- So, mixture is nice!


## Horse Race

- Horses: $1,2, \ldots, m$


[^0]
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- Odds: $o_{1}, o_{2}, \ldots, o_{m}$


[^1]
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- KT strategy: $q_{\mathrm{KT}}\left(y^{T}\right):=\int_{\Delta_{m-1}} p_{\boldsymbol{\theta}}\left(y^{T}\right) \mathrm{d} \pi(\boldsymbol{\theta})$, where $\pi(\boldsymbol{\theta})=\operatorname{Dir}\left(\boldsymbol{\theta} \left\lvert\, \frac{1}{2}\right., \ldots, \frac{1}{2}\right)$

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## Stock Investment

- Stocks: $1,2, \ldots, m$

Selected asset performance since Jan 3 high for S\&P 500


Image credit: https://www.reuters.com/article/usa-stocks-bearmarket-idCAKCN2N61PI

## Stock Investment

- Stocks: $1,2, \ldots, m$
- Price relatives (market vector):

$$
\begin{aligned}
\mathbf{x}_{t} & =\left(x_{t 1}, \ldots, x_{t m}\right) \in \mathcal{M} \subseteq \mathbb{R}_{\geq 0}^{m} \\
x_{t i} & :=\frac{(\text { end price of stock } i \text { on day } t)}{(\text { start price of stock } i \text { on day } t)}
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- Cumulative wealth: starting with $\$ \mathrm{~W}_{0}$,

$$
\mathrm{W}\left(\mathrm{x}^{T}\right)=\mathrm{W}_{0} \prod_{t=1}^{T}\left\langle\mathrm{~b}\left(\mathrm{x}^{t-1}\right), \mathrm{x}_{t}\right\rangle
$$

## Special Cases



## From Probability Assignment to Portfolio Selection

- By distributive law,

$$
\mathrm{W}\left(\mathbf{x}^{T}\right)=\mathrm{W}_{0} \prod_{t=1}^{T}\left\langle\mathrm{~b}\left(\mathrm{x}^{t-1}\right), \mathrm{x}_{t}\right\rangle=\mathrm{W}_{0} \sum_{y^{T} \in[m]^{T}}\left(\prod_{t=1}^{T} b\left(y_{t} \mid \mathrm{x}^{t-1}\right)\right) \mathbf{x}^{T}\left(y^{T}\right)
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where $\mathbf{x}^{T}\left(y^{T}\right):=x_{1 y_{1}} \ldots x_{T y_{T}}=\left(\right.$ multiplicative gain of the extremal portfolio $\left.y^{T}\right)$

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- A probability induced portfolio: for a probability $q\left(y^{T}\right)$, define

$$
\mathrm{W}^{q}\left(\mathbf{x}^{T}\right):=\mathrm{W}_{0} \sum_{y^{T} \in[m]^{T}} q\left(y^{T}\right) \mathbf{x}^{T}\left(y^{T}\right),
$$

which is achieved by a causal bettor $\mathrm{b}^{q}$ defined to satisfy

$$
\mathbf{W}^{q}\left(\mathbf{x}^{t}\right)=\mathbf{W}^{q}\left(\mathbf{x}^{t-1}\right)\left\langle\mathbf{b}^{q}\left(\mathbf{x}^{t-1}\right), \mathbf{x}_{t}\right\rangle
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## Portfolio Selection $\equiv$ Probability Assignment

Theorem

$$
\sup _{p \in \mathcal{P}} \sup _{\mathbf{x}^{T}} \frac{\mathrm{~W}^{p}\left(\mathbf{x}^{T}\right)}{\mathrm{W}^{q}\left(\mathbf{x}^{T}\right)}=\sup _{p \in \mathcal{P}} \sup _{y^{T}} \frac{p\left(y^{T}\right)}{q\left(y^{T}\right)}
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## Proof

$$
\begin{aligned}
& \sup _{\mathbf{x}^{n}} \sup _{p \in \mathcal{P}} \frac{\mathrm{~W}^{p}\left(\mathbf{x}^{n}\right)}{\mathrm{W}^{q}\left(\mathbf{x}^{n}\right)} \geq \sup _{y^{n} \in[m]^{n}} \sup _{p \in \mathcal{P}} \frac{\mathrm{~W}^{p}\left(\mathbf{e}_{y_{1}} \ldots \mathbf{e}_{y_{n}}\right)}{\mathrm{W}^{q}\left(\mathbf{e}_{y_{1}} \ldots \mathbf{e}_{y_{n}}\right)}=\sup _{y^{n} \in[m]^{n}} \sup _{p \in \mathcal{P}} \frac{p\left(y^{n}\right)}{q\left(y^{n}\right)} \\
& \sup _{\mathbf{x}^{n}} \sup _{p \in \mathcal{P}} \frac{\mathrm{~W}^{p}\left(\mathbf{x}^{n}\right)}{\mathrm{W}^{q}\left(\mathbf{x}^{n}\right)}=\sup _{\mathbf{x}^{n}} \sup _{p \in \mathcal{P}} \frac{\sum_{y^{n}} p\left(y^{n}\right) \mathbf{x}\left(y^{n}\right)}{\sum_{y^{n}} q\left(y^{n}\right) \mathbf{x}\left(y^{n}\right)} \leq \sup _{p \in \mathcal{P}} \sup _{y^{n}} \frac{p\left(y^{n}\right)}{q\left(y^{n}\right)}
\end{aligned}
$$

## Lemma * (Cover, 2006, Lemma 16.7.1)

For $a_{i}, b_{i} \geq 0$, we have $\frac{\sum_{i=1}^{n} a_{i}}{\sum_{i=1}^{n} b_{i}} \leq \max _{j \in[n] \frac{a}{b_{j}}}^{b_{j}}$, where $\frac{0}{0}:=0$

## Example: Constant Rebalanced Portfolios

- $\mathcal{P}_{\text {i.i.d. }}=\{$ i.i.d. categorical probabilities $\}=\left\{p_{\boldsymbol{\theta}}(\cdot): \boldsymbol{\theta} \in \Delta_{m-1}\right\}$


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- Cover's universal portfolio (Cover, 1991; Cover and Ordentlich, 1996): $\mathbf{b}^{\mathrm{UP}}:=\mathbf{b}^{q_{\kappa \mathrm{T}}}$

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- Time complexity: $O\left(t^{m-1}\right)$ at round $t$


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- Time complexity: $O\left(t^{m-1}\right)$ at round $t$
- Note: for horse race, UP is equivalent to the simple KT strategy


## From Universal Gambling to Confidence Sequences

## Supermartingales from Gambling

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## Proof.

For every $t, \mathrm{E}\left[\mathrm{W}_{t} \mid \mathbf{x}^{t-1}\right]=\mathrm{W}_{t-1}\left\langle\mathbf{b}_{t}, \mathrm{E}\left[\mathbf{x}_{t} \mid \mathbf{x}^{t-1}\right]\right\rangle \leq \mathrm{W}_{t-1}\left\langle\mathbf{b}_{t}, \mathbb{1}\right\rangle=\mathrm{W}_{t-1}$

## Examples

- Coin betting: $\mathbf{x}_{t}=\left(2 Y_{t}, 2\left(1-Y_{t}\right)\right), Y_{t} \in\{0,1\}$
- fair if $\mathrm{E}\left[Y_{t} \mid Y^{t-1}\right]=\frac{1}{2}$ (e.g., $Y_{t} \sim$ i.i.d. $\left.\operatorname{Bern}\left(\frac{1}{2}\right)\right)$


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- If $m=\mu$, any wealth process from $\operatorname{CTHR}(m)$ is martingale
- If $m \neq \mu$, there exists a causal betting strategy whose wealth process from $\operatorname{CTHR}(m)$ is strictly submartingale


## High-Level Intuition (Waudby-Smith and Ramdas, 2020b)

- For $\operatorname{CTHR}(m)$, we play a strategy $\left(\mathbf{b}\left(Y^{t-1} ; m\right)\right)_{t=1}^{\infty}$ and get $\left(\mathrm{W}\left(Y^{t} ; m\right)\right)_{t=1}^{\infty}$


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- If we collect all $m$ whose corresponding wealth never exceeds $\mathrm{W}_{0} / \delta$ by then, it forms a time-uniform confidence set with level $1-\delta$


## Confidence Sequence from CTHR $(m)$

- Formally, if we define

$$
C_{t}\left(Y^{t} ; \delta\right):=\left\{m \in(0,1): \sup _{1 \leq i \leq t} \frac{\mathrm{~W}\left(\mathbf{x}^{i} ; m\right)}{\mathrm{W}_{0}}<\frac{1}{\delta}\right\},
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- Orabona and Jun (2021) empirically showed that applying Cover's UP gives tight confidence sequences


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## Proof.

- Apply Ville's inequality
- The set is an interval, since $m \mapsto \phi_{t}\left(x ; \frac{1}{m}, \frac{1}{1-m}\right)$ is log-convex
${ }^{1}$ The optimal order is $\frac{1}{t} \log \log t$, which is implied by the law of iterated logarithm (LIL)


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- Note: the size of the interval behaves as $\sqrt{\frac{2}{t} \log \frac{1}{\delta}+\frac{1}{t} \log t+o(1)}$ for $t \gg 1$, which is comparable to $\sqrt{\frac{2}{t} \log \frac{1}{\delta}}$ from the standard Hoeffding ${ }^{1}$

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- Recall that Cover's UP is defined as a mixture of wealths of CRPs
- Consider a tight lower bound of the CRP wealth and take a mixture over the lower bounds

A Lower Bound on the Wealth of CRP

- Let $\bar{a}:=1-a$ for any $a \in \mathbb{R}$


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## Lemma (Generalization of (Waudby-Smith and Ramdas, 2020b, Lemma 1))

For any $n \in \mathbb{N}$ and $m \in(0,1)$, we have
$\log \left(b \frac{y}{m}+\bar{b} \frac{\bar{y}}{\bar{m}}\right) \geq \sum_{k=1}^{2 n-1} \frac{1}{k}\left(1-\frac{\bar{b}}{\bar{m}}\right)^{k}\left\{\left(1-\frac{y}{m}\right)^{2 n}-\left(1-\frac{y}{m}\right)^{k}\right\}+\left(1-\frac{y}{m}\right)^{2 n} \log \frac{\bar{b}}{\bar{m}}$
if $b \in[m, 1)$ and $y \geq 0$

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- Lower-bound the logarithm by moments of $y$, i.e., $\left(1, y, \ldots, y^{2 n}\right)$


## Key Lemma for the Proof

## Lemma (Generalization of (Fan et al., 2015, Lemma 4.1))

For an integer $\ell \geq 1$, if we define

$$
f_{\ell}(t):= \begin{cases}\left(\log (1+t)-\sum_{k=1}^{\ell-1}(-1)^{k+1} \frac{t^{k}}{k}\right) /\left((-1)^{\ell} \frac{t^{\ell}}{\ell}\right) & \text { if } t>-1 \text { and } t \neq 0 \\ -1 & \text { if } t=0\end{cases}
$$

then $t \mapsto f_{\ell}(t)$ is continuous and strictly increasing over $(-1, \infty)$

- Note: Fan et al. (2015) considered $\ell=2$, i.e.,

$$
f_{2}(t)= \begin{cases}\frac{\log (1+t)-t}{t^{2} / 2} & \text { if } t>-1 \text { and } t \neq 0 \\ -1 & \text { if } t=0\end{cases}
$$

## A Lower Bound on the Cumulative Wealth of CRP

- Since it is easy to check $\phi_{n}(x ; \boldsymbol{\rho}, \eta) \phi_{n}\left(x ; \boldsymbol{\rho}^{\prime}, \eta^{\prime}\right)=\phi_{n}\left(x ; \boldsymbol{\rho}+\boldsymbol{\rho}, \eta+\eta^{\prime}\right)$,


## Lemma

For any $n \in \mathbb{N}, m \in(0,1), b \in[0,1]$, and $y^{t} \in[0,1]^{t}$, we have

$$
\log \frac{\mathbf{W}_{t}^{b}\left(y^{t} ; m\right)}{\mathbf{W}_{0}} \geq \log \phi_{n}\left(\frac{\bar{b}}{\bar{m}} ; \rho_{n}\left(y^{t} ; m\right), \eta_{n}\left(y^{t} ; m\right)\right)
$$

if $m<b<1$, where $\eta_{n}\left(y^{t} ; m\right):=\sum_{i=1}^{t}\left(1-\frac{y_{i}}{m}\right)^{2 n}$ and

$$
\left(\rho_{n}\left(y^{t} ; m\right)\right)_{k}:=\sum_{i=1}^{t}\left\{\left(1-\frac{y_{i}}{m}\right)^{2 n}-\left(1-\frac{y_{i}}{m}\right)^{k}\right\} \quad \text { for } k=1, \ldots, 2 n-1
$$

- Lower-bound the logarithm by moments of $y^{t}$, i.e., $\left(\sum_{i=1}^{t} y_{i}^{j}\right)_{j=1}^{2 n}$
- Complexity from $O(t)$ to $O(n)$


## A Mixture of Lower Bounds Approach

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\bar{m} Z_{n}\left(\rho_{n}\left(y^{t} ; m\right), \eta_{n}\left(y^{t} ; m\right)\right)+m Z_{n}\left(\rho_{n}\left(\bar{y}^{t} ; \bar{m}\right), \eta_{n}\left(\bar{y}^{t} ; \bar{m}\right)\right),
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- We can construct a time-uniform confidence interval using this "mixture of wealth lower bounds"!
- We call this $\operatorname{LBUP}(n)$, where $n$ is the approximation order


## Caveats

- Computational bottleneck: computing the normalization constant $Z_{n}(\rho, \eta)$ of the form

$$
\int_{0}^{1} x^{\eta} \exp \left(\sum_{k=0}^{2 n-1} a_{k} x^{k}\right) \mathrm{d} x
$$

- Hence, $O(1)$ per round in principle, but may take longer than running exact UP due to numerical integration steps
- Larger $n$ leads to better approximation, but with increased numerical instability; $n=2$ or $n=3$ empirically work well
- Bad approximation in a small sample regime
- Hybrid UP: run UP for the first few samples and switch to LBUP


## Evolution of Wealth Processes

- The horizontal lines indicate an example threshold $\ln \frac{1}{\delta} \approx 2.996$ for $\delta=0.05$


Figure: An i.i.d. Bern(0.25) process

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Figure: An i.i.d. Beta $(10,30)$ process

## Experiments

- Confidence sequences with level 0.95 (i.e., $\delta=0.05$ )
- CB: betting strategy from another gambling construction
- HR: KT strategy
- UP: exact Cover's UP strategy
- LBUP: proposed lower-bound approach
- HybridUP: run exact UP for the first few steps and switch to LBUP
- PRECiSE (Orabona and Jun, 2021)


## Experiments



Figure: With i.i.d. Bern(0.25) processes

## Experiments



Figure: With i.i.d. Beta(1,3) processes

## Experiments



Figure: With i.i.d. Beta $(10,30)$ processes

## Take-Home Messages

- Confidence sequence is an important tool in modern data science
- Gambling with respect to probability induced strategies $\equiv$ probability assignment
- Confidence sequences from universal portfolios are very tight with small samples, but suffers high complexity
- They can be "efficiently" approximated by a mixture of lower bounds approach!


## Take-Home Messages

- Confidence sequence is an important tool in modern data science
- Gambling with respect to probability induced strategies $\equiv$ probability assignment
- Confidence sequences from universal portfolios are very tight with small samples, but suffers high complexity
- They can be "efficiently" approximated by a mixture of lower bounds approach!
Q. Can we construct a time-uniform confidence set for bounded vectors? Yes!
Q. Can there be a gambling other than $\operatorname{CTHR}(m)$ that corresponds to some other statistics applications?


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[^2]:    ${ }^{1}$ The optimal order is $\frac{1}{t} \log \log t$, which is implied by the law of iterated logarithm (LIL)

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