# Time-Uniform Confidence Sequences from Universal Gambling

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January 6, 2023

Joint work with Alankrita Bhatt (Caltech)

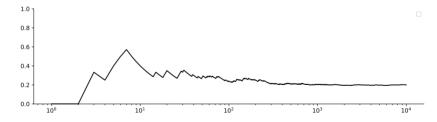
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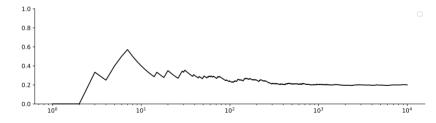
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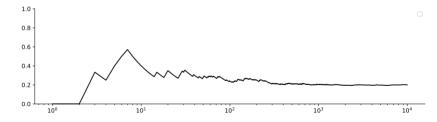


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- For reliable inference, we need to quantify confidence

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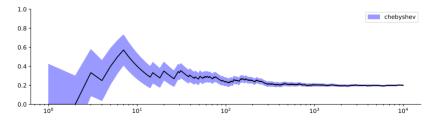
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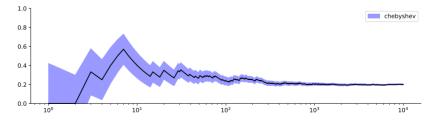
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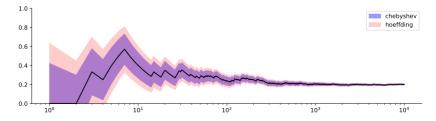
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- For such online data processing, we need to construct a sequence of confidence intervals that is valid at any time

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 Originally studied by Darling and Robbins (1967); Lai (1976), and recently resurrected by some statisticians (Ramdas et al., 2020; Waudby-Smith and Ramdas, 2020a,b; Howard et al., 2021) and computer scientists (Jun and Orabona, 2019; Orabona and Jun, 2021)

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#### Ville's inequality (Ville, 1939)

For a nonnegative supermartingale sequence  $(W_t)_{t=0}^{\infty}$  with  $W_0 > 0$ ,

$$\mathsf{P}\Big\{\sup_{t\geq 1}\frac{W_t}{W_0}\geq \frac{1}{\delta}\Big\}\leq \delta$$

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- Wait, what is gambling?
- As a slight detour, let's review canonical gambling problems and some universal gambling strategies

# Universal Gambling

## Coin Betting

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• Cumulative wealth: starting with \$W<sub>0</sub>,

$$\mathbf{W}_T = \mathbf{W}_0 \prod_{t=1}^T 2q(y_t | y^{t-1}) = \mathbf{W}_0 2^T q(y^T),$$

where  $q(y^T) := \prod_{t=1}^T q(y_t | y^{t-1})$ 



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• The best strategy is called minimax optimal

$$\min_{q} \max_{p \in \mathcal{P}} \max_{y^T} \log \frac{\mathsf{W}^p(y^T)}{\mathsf{W}^q(y^T)}$$

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 $\therefore$  universal compression  $\rightarrow$  universal betting!

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Asymptotically minimax optimal (Xie and Barron, 2000)

$$\max_{\theta \in [0,1]} \max_{y^T} \log \frac{p_{\theta}(y^T)}{q_{\mathsf{KT}}(y^T)} = \frac{1}{2} \log T + \frac{1}{2} \log \frac{\pi}{2} + o(1)$$

## Mixture Probability

• The KT probability  $q_{\mathrm{KT}}(\cdot|y^{t-1})$  is induced by a mixture probability, i.e.,

$$q_{\mathrm{KT}}(y^T) \equiv \int_0^1 p_{\theta}(y^T) \,\mathrm{d}\pi(\theta)$$

for  $\pi(\theta) = \text{Beta}(\theta|\frac{1}{2}, \frac{1}{2})$ 

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• The KT probability  $q_{\rm KT}(\cdot|y^{t-1})$  is induced by a mixture probability, i.e.,

$$q_{\mathrm{KT}}(y^T) \equiv \int_0^1 p_{\theta}(y^T) \,\mathrm{d}\pi(\theta)$$

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• So, mixture is nice!

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Selected asset performance since Jan 3 high for S&P 500



Note: Data as of June 13 morning trading Source: Refinitly

Image credit: https://www.reuters.com/article/usa-stocks-bearmarket-idCAKCN2N61PI

- Stocks:  $1, 2, \ldots, m$
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$$\mathbf{x}_t = (x_{t1}, \dots, x_{tm}) \in \mathcal{M} \subseteq \mathbb{R}^m_{\geq 0},$$
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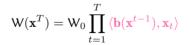
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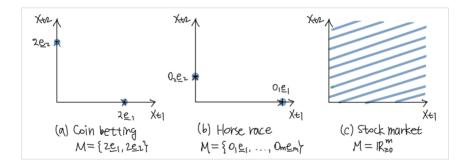


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## From Probability Assignment to Portfolio Selection

• By distributive law,

$$\mathsf{W}(\mathbf{x}^{T}) = \mathsf{W}_{0} \prod_{t=1}^{T} \langle \mathbf{b}(\mathbf{x}^{t-1}), \mathbf{x}_{t} \rangle = \mathsf{W}_{0} \sum_{y^{T} \in [m]^{T}} \left( \prod_{t=1}^{T} b(y_{t} | \mathbf{x}^{t-1}) \right) \mathbf{x}^{T}(y^{T}),$$

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• A probability induced portfolio: for a probability  $q(y^T)$ , define

$$\mathsf{W}^q(\mathbf{x}^T) := \mathsf{W}_0 \sum_{y^T \in [m]^T} q(y^T) \mathbf{x}^T(y^T),$$

which is achieved by a causal bettor  $\mathbf{b}^q$  defined to satisfy

$$\mathsf{W}^{q}(\mathbf{x}^{t}) = \mathsf{W}^{q}(\mathbf{x}^{t-1}) \langle \mathbf{b}^{q}(\mathbf{x}^{t-1}), \mathbf{x}_{t} \rangle$$

## Portfolio Selection $\equiv$ Probability Assignment

#### Theorem

$$\sup_{p \in \mathcal{P}} \sup_{\mathbf{x}^T} \frac{\mathsf{W}^p(\mathbf{x}^T)}{\mathsf{W}^q(\mathbf{x}^T)} = \sup_{p \in \mathcal{P}} \sup_{y^T} \frac{p(y^T)}{q(y^T)}$$

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Proof

$$\begin{split} \sup_{\mathbf{x}^n} \sup_{p \in \mathcal{P}} \frac{\mathsf{W}^p(\mathbf{x}^n)}{\mathsf{W}^q(\mathbf{x}^n)} &\geq \sup_{y^n \in [m]^n} \sup_{p \in \mathcal{P}} \frac{\mathsf{W}^p(\mathbf{e}_{y_1} \dots \mathbf{e}_{y_n})}{\mathsf{W}^q(\mathbf{e}_{y_1} \dots \mathbf{e}_{y_n})} = \sup_{y^n \in [m]^n} \sup_{p \in \mathcal{P}} \frac{p(y^n)}{q(y^n)} \\ \sup_{\mathbf{x}^n} \sup_{p \in \mathcal{P}} \frac{\mathsf{W}^p(\mathbf{x}^n)}{\mathsf{W}^q(\mathbf{x}^n)} &= \sup_{\mathbf{x}^n} \sup_{p \in \mathcal{P}} \frac{\sum_{y^n} p(y^n) \mathbf{x}(y^n)}{\sum_{y^n} q(y^n) \mathbf{x}(y^n)} \stackrel{(\star)}{\leq} \sup_{p \in \mathcal{P}} \sup_{y^n} \frac{p(y^n)}{q(y^n)} \end{split}$$

Lemma \* (Cover, 2006, Lemma 16.7.1)  
For 
$$a_i, b_i \ge 0$$
, we have  $\frac{\sum_{i=1}^n a_i}{\sum_{i=1}^n b_i} \le \max_{j \in [n]} \frac{a_j}{b_j}$ , where  $\frac{0}{0} := 0$ 

Jon Ryu

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- Note: for horse race, UP is equivalent to the simple KT strategy

# From Universal Gambling to Confidence Sequences

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#### Proof.

For every t,  $\mathsf{E}[\mathsf{W}_t|\mathbf{x}^{t-1}] = \mathsf{W}_{t-1}\langle \mathbf{b}_t, \mathsf{E}[\mathbf{x}_t|\mathbf{x}^{t-1}] \rangle \leq \mathsf{W}_{t-1}\langle \mathbf{b}_t, \mathbb{1} \rangle = \mathsf{W}_{t-1}$ 

- Coin betting:  $\mathbf{x}_t = (2Y_t, 2(1 Y_t)), Y_t \in \{0, 1\}$ 
  - fair if  $\mathsf{E}[Y_t|Y^{t-1}] = \frac{1}{2}$  (e.g.,  $Y_t \sim \text{i.i.d. Bern}(\frac{1}{2})$ )

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- Two-horse race:  $\mathbf{x}_t = (o_1 Y_t, o_2(1 Y_t)), Y_t \in \{0, 1\}$ • fair if  $\frac{1}{o_1} + \frac{1}{o_2} = 1$  and  $\mathsf{E}[Y_t|Y^{t-1}] = \frac{1}{o_1}$  (e.g.,  $Y_t \sim \text{i.i.d. Bern}(\frac{1}{o_1})$ )

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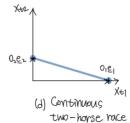
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- Continuous two-horse race:  $\mathbf{x}_t = (o_1 Y_t, o_2(1 Y_t)), Y_t \in [0, 1]$ 
  - fair if  $\frac{1}{o_1} + \frac{1}{o_2} = 1$  and  $E[Y_t|Y^{t-1}] = \frac{1}{o_1}$ ;
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• Recall: Assume  $\mathsf{E}[Y_t|Y^{t-1}] \equiv \mu$  for some  $\mu \in (0,1)$ 

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## High-Level Intuition (Waudby-Smith and Ramdas, 2020b)

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• If we collect all m whose corresponding wealth never exceeds  $W_0/\delta$  by then, it forms a time-uniform confidence set with level  $1 - \delta$ 

• Formally, if we define

$$C_t(Y^t;\delta) := \Big\{ m \in (0,1) \colon \sup_{1 \le i \le t} \frac{\mathsf{W}(\mathbf{x}^i;m)}{\mathsf{W}_0} < \frac{1}{\delta} \Big\},$$

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## Confidence Sequence from KT Betting

#### Theorem

 $(C^{\mathrm{KT}}_t(Y^t;\delta))_{t=1}^\infty$  is a time-uniform confidence interval with level  $1-\delta$ 

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- Note: the size of the interval behaves as  $\sqrt{\frac{2}{t}\log\frac{1}{\delta} + \frac{1}{t}\log t + o(1)}$  for  $t \gg 1$ , which is comparable to  $\sqrt{\frac{2}{t}\log\frac{1}{\delta}}$  from the standard Hoeffding<sup>1</sup>

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- An alternative approach (Ryu and Bhatt, 2022)
  - Recall that Cover's UP is defined as a mixture of wealths of CRPs
  - Consider a tight lower bound of the CRP wealth and take a mixture over the lower bounds

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- Lower-bound the logarithm by moments of y, i.e.,  $(1, y, \dots, y^{2n})$

Jon Ryu

#### Key Lemma for the Proof

Lemma (Generalization of (Fan et al., 2015, Lemma 4.1)) For an integer  $\ell \ge 1$ , if we define

$$f_{\ell}(t) := \begin{cases} \left(\log(1+t) - \sum_{k=1}^{\ell-1} (-1)^{k+1} \frac{t^k}{k}\right) \middle/ \left((-1)^{\ell} \frac{t^{\ell}}{\ell}\right) & \text{if } t > -1 \text{ and } t \neq 0, \\ -1 & \text{if } t = 0, \end{cases}$$

then  $t\mapsto f_\ell(t)$  is continuous and strictly increasing over  $(-1,\infty)$ 

• Note: Fan et al. (2015) considered  $\ell = 2$ , i.e.,

$$f_2(t) = \begin{cases} \frac{\log(1+t) - t}{t^2/2} & \text{if } t > -1 \text{ and } t \neq 0, \\ -1 & \text{if } t = 0 \end{cases}$$

## A Lower Bound on the Cumulative Wealth of CRP

• Since it is easy to check  $\phi_n(x; \rho, \eta)\phi_n(x; \rho', \eta') = \phi_n(x; \rho + \rho, \eta + \eta')$ ,

#### Lemma

For any  $n \in \mathbb{N}, m \in (0,1), b \in [0,1],$  and  $y^t \in [0,1]^t,$  we have

$$\log \frac{\mathsf{W}_t^b(y^t;m)}{\mathsf{W}_0} \geq \log \phi_n\Big(\frac{\bar{b}}{\bar{m}}; \boldsymbol{\rho}_n(y^t;m), \eta_n(y^t;m)\Big)$$

if m < b < 1, where  $\eta_n(y^t;m) := \sum_{i=1}^t {(1-\frac{y_i}{m})^{2n}}$  and

$$(\boldsymbol{\rho}_n(y^t;m))_k := \sum_{i=1}^t \left\{ \left(1 - \frac{y_i}{m}\right)^{2n} - \left(1 - \frac{y_i}{m}\right)^k \right\} \text{ for } k = 1, \dots, 2n-1$$

- Lower-bound the logarithm by moments of  $y^t$ , i.e.,  $(\sum_{i=1}^t y_i^j)_{j=1}^{2n}$
- Complexity from O(t) to O(n)

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- We call this LBUP(*n*), where *n* is the approximation order

#### Caveats

• Computational bottleneck: computing the normalization constant  $Z_n(\rho, \eta)$  of the form

$$\int_0^1 x^\eta \exp\left(\sum_{k=0}^{2n-1} a_k x^k\right) \mathrm{d}x$$

- Hence, O(1) per round in principle, but may take longer than running exact UP due to numerical integration steps
- Larger n leads to better approximation, but with increased numerical instability; n = 2 or n = 3 empirically work well
- Bad approximation in a small sample regime
  - Hybrid UP: run UP for the first few samples and switch to LBUP

#### **Evolution of Wealth Processes**

• The horizontal lines indicate an example threshold  $\ln \frac{1}{\delta} \approx 2.996$  for  $\delta = 0.05$ 

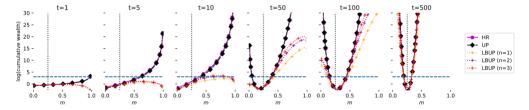


Figure: An i.i.d. Bern(0.25) process

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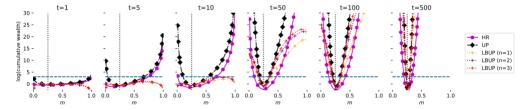


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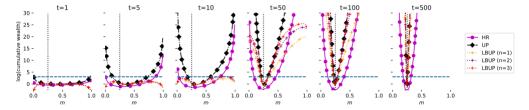


Figure: An i.i.d. Beta(10,30) process

- Confidence sequences with level 0.95 (i.e.,  $\delta=0.05)$
- CB: betting strategy from another gambling construction
- HR: KT strategy
- UP: exact Cover's UP strategy
- LBUP: proposed lower-bound approach
- HybridUP: run exact UP for the first few steps and switch to LBUP
- PRECiSE (Orabona and Jun, 2021)

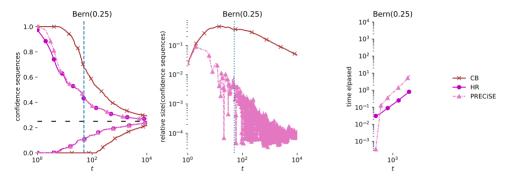


Figure: With i.i.d. Bern(0.25) processes

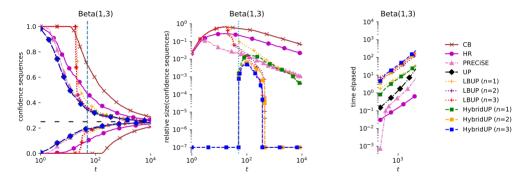


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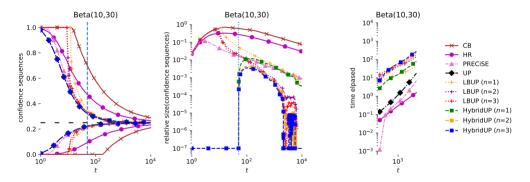


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### Take-Home Messages

- Confidence sequence is an important tool in modern data science
- Gambling with respect to probability induced strategies  $\equiv$  probability assignment
- Confidence sequences from universal portfolios are very tight with small samples, but suffers high complexity
- They can be "efficiently" approximated by a mixture of lower bounds approach!

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- Q. Can we construct a time-uniform confidence set for bounded vectors? Yes!
- **Q.** Can there be a gambling other than CTHR(m) that corresponds to some other statistics applications?

## References I

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