Parameter-Free Online Linear Optimization

with Side Information via Universal Coin Betting

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PROBLEM: ONLINE LINEAR OPTIMIZATION (OLO)

- Assume Hilbert space V with norm || · ||
- In each round t = 1,2,...
- Learner picks action $\mathbf{w}_t \in V$
- Receives a vector $\mathbf{g}_t \in V$ such that $||\mathbf{g}_t|| \le 1$
- Gains reward $\langle \mathbf{g}_t, \mathbf{w}_t \rangle$
- Goal: maximize the cumulative reward $\sum_{t=1}^{T} \langle \mathbf{g}_t, \mathbf{w}_t \rangle$
- The standard metric: regret with respect to the best static competitor in hindsight

$$\mathsf{Reg}(\mathbf{u}; \mathbf{g}^T) := \sum_{t=1}^T \langle \mathbf{g}_t, \mathbf{u}
angle - \sum_{t=1}^T \langle \mathbf{g}_t, \mathbf{w}_t
angle ext{ for } \mathbf{u} \in V$$

- Two issues
- 1. Learning rate tuning requires a priori knowledge on ||u||
- 2. Static competitors are weak

✓ Parameter-Free OLO via Universal Coin Betting

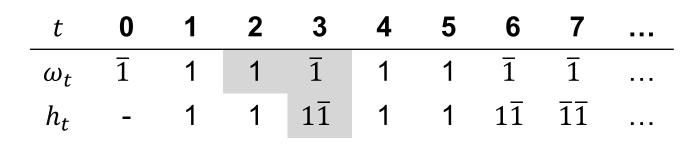
- To attain optimal rates, naive approaches require ||u||
- Q. Can we attain optimal regret w/o the need of tuning parameters?
- A. Orabona and Pál (2016) showed that a universal coin betting algorithm can be converted to a near-optimal-regret parameter-free OLO algorithm!
- Key tool: Fenchel duality
- Note: Other parameter-free algorithms exist

✓ OLO with Side Information

- Static competitors $\{\mathbf{u} : \mathbf{u} \in V\}$ are weak
 - **Example**: for g, -g, g, -g, ..., the best reward with $u \in V$ is zero
 - In general, $\langle \sum_{t=1}^T \mathbf{g}_t , \mathbf{u}_t \rangle$ can be large iff $||\sum_{t=1}^T \mathbf{g}_t ||$ is large
- Q. Can we leverage a possible structure in $(\mathbf{g}_t)_{t>1}$?
- Our approach: Provided that we have access to a side information sequence $(h_t \in \{1, \overline{1}\})_{t \ge 1}$ which may potentially capture a structure, develop a method that adapts to side information!
- Example: $h_t = \operatorname{sgn}(\langle \mathbf{g}_{t-1}, \mathbf{f} \rangle)$ (quantization with $\mathbf{f} \in V$)
- To capture a more complex structure, we consider:
- **Def (tree side information):** Given a suffix tree **T** and an auxiliary sequence $\Omega = (\omega_t \in \{1, \overline{1}\})_{t \ge 1}$, the tree side information $(h_t)_{t \ge 1}$ is $h_t = (\text{the matching suffix of } \omega_{-\infty}^t \text{ w.r.t. } \mathbf{T})$

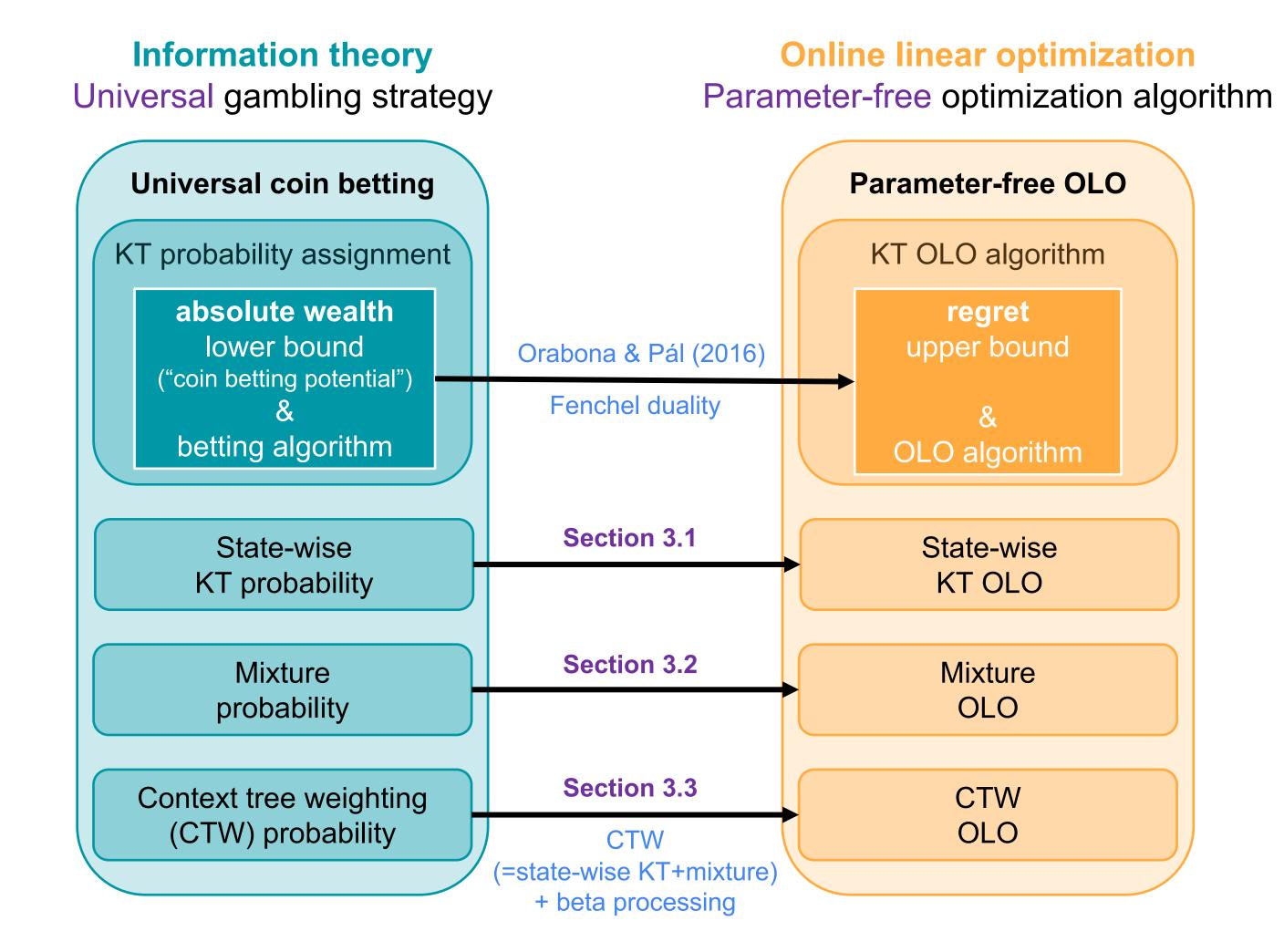
 n_t –

• Example:



$\begin{array}{c} *1 \\ \hline 1\overline{1} \\ \hline 1\overline{1} \end{array}$ $\mathbf{T} = \{1\overline{1}, \overline{1}\overline{1}, * 1\}$

PARAMETER-FREE OLO WITH SIDE INFORMATION VIA UNIVERSAL COIN BETTING



- Idea: coin betting wealth lower bound can be translated into OLO algorithm AND regret bound
- In general, parameter-freeness incurs additional multiplicative logarithmic factors
- Building blocks
 - To adapt to a single side information sequence: state-wise KT OLO
 - To adapt to any one of multiple side information sequences: mixture OLO
- Application: tree side information
 - **Goal**: given Ω , adapt to any tree side information sequence of depth $\leq D$
 - Approach: Take a mixture of state-wise KT OLOs for all tree side information sequences, following CTW (Willems et al. 1995) for universal tree source compression
 - Challenge: The mixture over all subtrees of depth $\leq D$ involves $O(2^{2^D})$ summands
 - CTW OLO algorithm
 - → can adapt the beta processing algorithm (Willems et al. 2006) for CTW;
 - \rightarrow runs in O(D) time complexity per step, with O(D) storage complexity



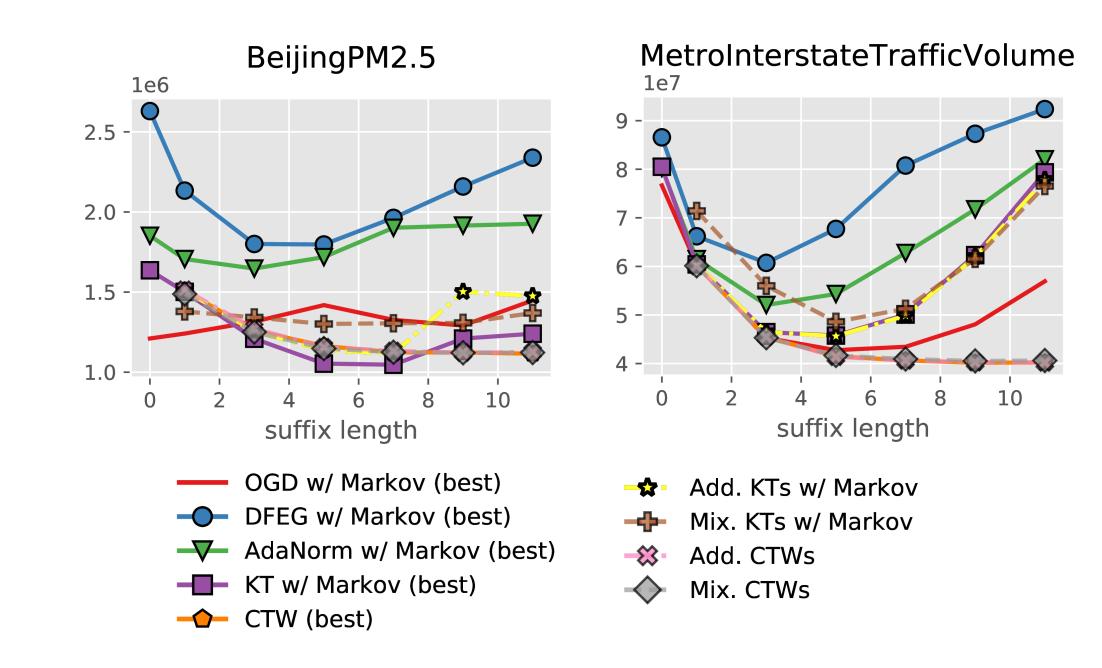


EXPERIMENT

• Online linear regression with absolute loss: $\mathbf{x} \in \mathbb{R}^d$, $y \in \mathbb{R}$

$$\ell_t(\mathbf{w}_t) \triangleq \ell(\hat{y}_t, y_t) \quad \hat{y}_t \triangleq \langle \mathbf{w}_t, \mathbf{x}_t \rangle$$
$$\partial \ell_t(\mathbf{w}_t) = \operatorname{sgn}(\langle \mathbf{w}_t, \mathbf{x}_t \rangle - y_t) \mathbf{x}_t$$

- Two real-world temporal datasets
- 1. Beijing PM2.5 (air pollution dataset)
- 2. Metro Inter State Traffic Volume (traffic volume dataset)
- Auxiliary sequence construction: for each dimension $i \in [d]$, apply canonical binary quantizer $Q_{\mathbf{e}_i}$ for each symbol (\mathbf{e}_i = the i-th standard vector)
- Run algorithms with tree side information of depth $D \in \{0,1,3,5,7,9,11\}$
- Since we do not know which depth is best a priori, apply mixture (or addition)



Observations

- 1. The performance of OGD, DFEG, AdaNorm with Markov side information get worse as the side information depth increases
- 2. The best of CTWs over dimensions achieves incurs almost the lowest losses
- 3. The addition or mixture of CTWs over the dimensions attain the performance of the best of optimally tuned OGDs, KTs, and CTWs

References

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